Within-Problem Learning for Efficient Lower Bound Computation in Max-SAT Solving

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Abstract
This paper focuses on improving branch-and-bound Max-SAT solvers by speeding up the lower bound computation. We notice that the existing propagation-based computing methods and the resolution-based computing methods, which have been studied intensively, both suffer from several drawbacks. In order to overcome these drawbacks, we propose a new method with a nice property that guarantees the increment of lower bounds. The new method exploits within-problem learning techniques. More specifically, at each branch point in the search-tree, the current node is enabled to inherit inconsistencies from its parent and learn information about the effectiveness of the lower bound computing procedure from previous nodes. Furthermore, after branching on a new variable, the inconsistencies may shrink by applying unit propagation to them, and such process increases the probability of getting better lower bounds. We graft the new techniques into maxsatz and the experimental results demonstrate that the new solver outperforms the best state-of-the-art solvers on a wide range of instances including random and structured ones.

Introduction
In recent years there has been an increasing interest in studying the Max-SAT problem. On one hand, more new applications of Max-SAT solvers have been found to areas as diverse as machine learning (Yang, Wu, & Jiang 2007), circuit debugging (Safarpour et al. 2007), and bioinformatics (Zhang et al. 2006). On the other hand, remarkable progress has been achieved in designing and implementing efficient exact Max-SAT solvers (Xing & Zhang 2005; Larrosa & Heras 2006; Larrosa, Heras, & de Givry 2008; Heras, Larrosa, & Oliveras 2008; Lin & Su 2007; Li, Manyà, & Planes 2006; 2007; Darras et al. 2007; Pipatsrisawat & Darwiche 2007). The Max-SAT Evaluation 2006† and 2007‡ advanced and witnessed such progress.

The most common approach for solving Max-SAT is based on a branch and bound (BnB) algorithm, which is a general scheme for NP-hard optimization problems. Results from the Max-SAT Evaluation indicate that BnB solvers are generally faster than other approaches. The key factor that makes BnB solvers efficient is to compute lower bounds of good quality. In (Li, Manyà, & Planes 2005; 2006), unit propagation and failed literal detection, called propagation-based methods, are exploited to compute lower bounds. These procedures are able to identify more disjoint inconsistent subformulas for a given formula, thus improve previous lower bounds significantly. However, they still have large overheads. First, due to the lack of mechanism to maintain these inconsistent subformulas, such inconsistencies may be rediscovered again and again by the descendant nodes. Second, at a search node, if the lower bound obtained is less than the upper bound, the lower bound computation effort is completely wasted.

Another method, called resolution-based method, to compute lower bound is using inference rules to simplify formulas and make the conflict explicit (Larrosa & Heras 2006; Li, Manyà, & Planes 2007). This method may avoid the above problems by applying rules to the inconsistent subformulas. In this case, the explicit empty clauses make the lower bound in the subsequent search incremental, since the inconsistent subformulas involved are not encountered again and the empty clauses are inherited by the descendant nodes. Unfortunately, this method is limited because the inference rules cannot be applied unless a formula matches some special patterns. Moreover, applying rules may be time and space consuming because more new clauses have to be added to the current formula and removed upon backtracking. So, although complete resolution has been proposed by (Bonet, Levy, & Manyà 2007), only practically efficient rules are implemented in existing solvers.

In this paper, we attempt to propose a new algorithm for Max-SAT, which aims to address all the above problems. For this purpose, the new lower bound computation method in our algorithm must satisfy at least three requirements:

†This work was partially supported by National Basic Research 973 Program of China under grant 2005CB321902, National Natural Science Foundation of China grants 60725207 and 60473004, the Australian Research Council grant DP0452628, Guangdong Provincial Natural Science Foundation grants 04205407 and 06023195, and Start-up research fund for induced talents of Jinan University.

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†http://www.iiia.csic.es/~maxsat06/
‡http://www.maxsat07.udl.es/
First, it computes lower bounds in an incremental way; second, it can be applied to all types of inconsistent subformulas; third, it can be implemented efficiently.

The basic idea is to utilize the so-called within-problem learning techniques, which mean, at the current node, exploiting information gathered by other nodes of the search tree to boost the whole depth-first search. This kind of techniques are widely used in SAT and CSP study. For example, clause learning, a well-known within-problem learning method, has been incorporated in the modern SAT (Zhang et al. 2001) and Max-SAT solvers (Argelich & Manya 2007). In our algorithm, each node of the search tree learns inconsistencies from its parent, and eliminates the clauses which are redundant in the current inconsistencies. Moreover, the failed literal detection may be time consuming. Thus, our algorithm also enables each node of the search tree to learn the information about what the effectiveness of the failed literal detection is for the already explored nodes, and decide whether such procedure should be invoked. We graft these new techniques into the Max-SAT solver maxsatz (Li, Manya, & Planes 2006; 2007), and the experimental results demonstrate that our new techniques enhance the effectiveness.

The paper is organized as follows. In the next section, we introduce some necessary background and related work. Then we present our within-problem learning algorithms formally. We have implemented new Max-SAT solvers called incmaxsatz by building on maxsatz. Experimental results demonstrating the performance of our solvers are presented next followed by some final conclusions.

**Preliminaries and Related Work**

In propositional logic, a variable \( x \) may take values true or false. A literal is either a variable \( x \) or its negation \( \bar{x} \). A clause \( C = l_1 \lor l_2 \lor \ldots \lor l_k \) is a disjunction of literals. A CNF formula is a conjunction of clauses, which is usually represented as a set of clauses \( \{C_1, C_2, \ldots, C_m\} \). For a formula \( F \) over a variable set \( V \), an assignment is a mapping from \( V \) to \( \{0, 1\} \). The assignment is complete if \( V' = V \); otherwise it is partial. An assignment satisfies a clause if it satisfies at least one literal in the clause, and satisfies a CNF formula if it satisfies all clauses. The instantiation of a formula \( F \) by forcing literal \( l \) to be true, denoted by \( F[l] \), produces a new formula by removing all clauses containing \( l \), and removing \( l \) from each clause where it occurs. This procedure is also called the application of one-literal rule. A formula is inconsistent if it cannot be satisfied by any assignment.

Given a CNF formula, the Max-SAT problem is to find an assignment that maximizes the number of satisfied clauses, or equivalently, minimizes the number of unsatisfied clauses.

The basic BnB algorithm for Max-SAT solving behaves as follows: Given a formula \( F \), BnB traverses the search tree representing the space of all possible assignments in a depth-first manner. At each node, BnB computes the lower bound \( LB \), which is the underestimation of the minimum number of unsatisfied clauses if the current partial assignment is extended to a complete one. Then \( LB \) is compared with the upper bound \( UB \), which is the minimum number of unsatisfied clauses found so far for a complete assignment. If \( LB \geq UB \), BnB backtracks to a higher level in the search tree; otherwise, an unassigned variable is chosen to assign a value and the search continues. After the entire space is explored, the value \( UB \) takes is the minimum number of unsatisfied clauses in \( F \).

As mentioned in the previous section, \( LB \) can be computed by exploiting unit propagation (UP) and failed literals detection (FL) to search for disjoint inconsistent subformulas. We briefly describe these two procedures as follows.

- **UP** (Li, Manya, & Planes 2005): Given a formula \( F \), UP maintains a set \( UNIT \), which contains the unit clauses derived so far, and repeatedly applies one-literal rule, for each unit clause in \( UNIT \), to simplify \( F \) until no more unit clause or an empty clause is derived. If an empty clause is derived, the set of clauses that used to derive this empty clause forms an inconsistent subformula \( IS \). In order to identify other disjoint inconsistent subformulas, \( IS \) is removed from \( F \), and the iterative process continues. If no more empty clause can be derived, then \( LB \) is equal to the number of empty clauses in \( F \) plus the number of inconsistent subformulas identified.

- **FL** (Li, Manya, & Planes 2006): Given a formula \( F \), FL applies unit propagation to \( F \cup \{x\} \) and \( F \cup \bar{x} \) for every variable \( x \) occurring in \( F \). If these two processes both derive an empty clause, then \( IS_x \cup IS_{\bar{x}} \setminus \{x, \bar{x}\} \) is an inconsistent subformula, where \( IS_x \) and \( IS_{\bar{x}} \) are the inconsistent subformulas detected in \( F \cup \{x\} \) and \( F \cup \bar{x} \) respectively.

Note that in (Li, Manya, & Planes 2006), FL is invoked after UP at every node unless UP leads to backtracking. In our algorithm, at each node, whether FL is executed depends on the effect FL had on the nodes which have been explored. Another new solver called maxsatz14iess (Darras et al. 2007), which memorizes inconsistencies containing at most 5 clauses and no branching variables, will be compared with our work in the following sections.

**Within-problem Learning Algorithms**

**Learning from One’s Own Parent**

In classical SAT solving, the solver utilizes unit propagation to simplify a formula, so the propagation procedure at each node of the search tree is completely different from that at its parent node. While in Max-SAT solving, unit propagation is exploited to compute bounds. Since such computation does not change the formula, the propagation procedure at each node is very close to its parent. Thus, computing lower bounds at each node is very close to its parent. Thus, computing lower bounds at each node from scratch will waste a lot of time to re-apply unit propagation to the same clauses. This also happens in failed literal detection. Moreover, applying unit propagation from scratch may cause another problem: The lower bound may decrease as the search-depth increases, i.e., as more variables are instantiated.

**Example 1** Let \( F = \{x_1, x_1 \lor x_2, x_1 \lor x_2 \lor x_3, x_3 \lor x_4, x_2 \lor x_4, x_5 \lor x_2, x_5 \lor x_4, x_5 \lor x_6, x_6 \lor x_7, x_6 \lor x_7 \lor x_8, x_8 \lor x_9, x_5 \lor x_8 \lor x_9\} \). At first, \( UNIT = \{x_1, x_5\} \).
$x_1$ is propagated first, and we get an inconsistent subformula $IS_1 = \{x_1, x_2, \lor x_2, \lor x_2, \lor x_2, x_3, \lor x_4, \lor x_5, \lor x_6\}$. Then $x_5$ is propagated, and another inconsistent subformula $IS_2 = \{x_3, \lor x_3, \lor x_4, \lor x_7, \lor x_8, \lor x_9, \lor x_9, \lor x_9\}$ is obtained. Suppose that we branch next on variable $x_1$. After assigning true to $x_1$, we get $F' = \{x_1, x_3, \lor x_4, \lor x_5, \lor x_6, \lor x_7, \lor x_8\}$, and the new unit clause $x_2$ is added to $UNIT$ such that $UNIT = \{x_3, x_2\}$. In the following, we show that the previous lower bound computation methods will decrease $LB$.

- **maxsat$^2$**: According to the unit clause selection heuristic in maxsat, $x_5$ will be propagated before $x_2$, because $x_5$ is “older” than $x_2$. After propagating $x_5$ on $F'$, new unit clauses $x_2$, $x_4$, and $x_6$ are added to $UNIT$, and then after propagating $x_2$ and $x_4$, we derive an empty clause and obtain an inconsistent subformula $IS' = \{x_2 \lor x_4, x_5 \lor x_2, x_5 \lor x_4\}$. After excluding $IS'$ from $F'$, the remaining clause-set is $\{x_2, x_3 \lor x_4, x_5 \lor x_6, x_7 \lor x_8, x_9 \lor x_5, x_9 \lor x_6, x_9 \lor x_6\}$, in which no addition inconsistency can be detected. Thus, the new lower bound $LB' = 1$, which is less than $LB$ obtained at its parent node.

- **maxsat$^{14}_{ics}$**: According to the inconsistent subformula storage heuristic in maxsat$^{14}_{ics}$, $IS_1$ is not stored because the branching variable $x_1$ occurs in it, and neither is $IS_2$ because the number of clauses in it exceeds five. The following processes are the same as in the maxsat$^2$ case.

As is known, lower bound computation plays an important role in reducing the size of the BnB search tree by making $LB$ increase quickly so as to equal or exceed $UB$. So it would be better to avoid, when the search-depth increases, decreasing $LB$ in lower bound computation. To the best of our knowledge, however, no previous method provides a mechanism to guarantee the increment of lower bounds, i.e., to guarantee that the lower bound at every node, except for the root, will be greater than or equal to the lower bound at its parent. In fact, there is a simple way to address this problem. At each node of the search-tree, we can memorize all the inconsistent subformulas after they are identified, such that each child-node is able to inherit all the inconsistencies of its parent first and then exploits unit propagation to find new inconsistencies. This process is safe because of the following observation.

**Proposition 1** If $F$ is an inconsistent formula, so are both $F[x]$ and $F[\overline{x}]$ for any variable $x$ in $F$.

This proposition trivially holds and guarantees that all inconsistencies of a node are still present in other nodes of the subtree rooted by this node. It is clear that memorizing the inconsistencies guarantees the increment of lower bounds.

**Example 2** In Example 1, after detecting $IS_1$ and $IS_2$, we memorize them. We still choose $x_1$, and assign true to it. Now $IS_2$ is left intact because it does not contain $x_1$, and $IS_1$ becomes $\{x_2, x_3 \lor x_2, x_3 \lor x_4, x_3 \lor x_4\}$ (by applying the one-literal rule). It is obvious that $IS_1$ and $IS_2$ are still inconsistent and $LB$ keeps the value 2.

Although the above method avoids re-computation and guarantees the increment of lower bound, it is still far from satisfactory because the inconsistent subformulas may contain redundant clauses, i.e., clauses whose removal will not make the subformula satisfiable.

**Example 3** Let $F$ be the formula in Example 1, and $IS_1$ and $IS_2$ detected and memorized as in the previous examples. Now suppose that we branch next on $x_7$ rather than $x_1$, and assign it false. Now $IS_1$ is left intact and $IS_2$ becomes $\{x_3, x_5 \lor x_6, x_6, x_8 \lor x_9, x_5 \lor x_8 \lor x_9\}$, in which the clauses $x_8 \lor x_9$ and $x_8 \lor x_9$ are redundant.

In the above example, if $IS_2$ can be reduced to $\{x_5, x_5 \lor x_6, x_6 \lor x_9\}$, the remaining clause-set is $\{x_5, x_5 \lor x_6, x_6 \lor x_9\}$, which may be used together with other clauses, to form new inconsistent subformulas such that $LB$ increases.

In order to reduce an inconsistent subformula, one may utilize the algorithms for extracting minimal unsatisfiable cores like in (Mneimneh et al. 2005). However, these techniques, for a BnB Max-SAT solver, is costly because the unsatisfiable cores should be extracted many times even at one search node. Fortunately, we may appeal to unit propagation again. Unit propagation also works here because detecting an inconsistent formula by unit propagation means proving its unsatisfiability by unit resolution$^4$, and we have the following proposition.

**Proposition 2** Let $F$ be an inconsistent formula. If the unsatisfiability of $F$ can be proved by unit resolution, so can $F[x]$ and $F[\overline{x}]$ for any variable $x$ in $F$.

Proposition 2 guarantees that the child-node can apply unit propagation to detect, hopefully smaller, inconsistent subformulas from the ones inherited from its parent.

**Example 4** In Example 3, applying unit propagation to the inconsistent subformula $\{x_3, x_5 \lor x_6, x_6 \lor x_9, x_5 \lor x_8 \lor x_9\}$ will identify a smaller one $\{x_3, x_5 \lor x_6, x_6\}$.

Apparently, the new lower bound computation method can avoid re-computation and guarantee the increment of lower bounds. Furthermore, thanks to the unit propagation to reduce inconsistencies, the memorized inconsistent subformulas may become smaller and smaller as more and more variables are instantiated. This means that we are more likely to get more inconsistent subformulas and improve lower bounds.

We found that, from the empirical study, it is not a good choice to memorize inconsistencies when $LB$ is much less than $UB$. An intuitive explanation for this phenomena is that “a child should not learn from his parent when his parent does not act as a good example”. So in our algorithm, the learning is triggered when $LB \leq UB \alpha$. The best value of $\alpha$ is instance-dependent, nevertheless we simply set $\alpha = 0.3$ and $\alpha = 0.8$ for Max-2-SAT and Max-3-SAT respectively. The empirical results show that such setting seems good enough.

$^4$Unit resolution states that from $l \lor f$, and $l$, where $l$ is a literal and $D$ is a disjunction of literals, we can derive $D$. 

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No efficient inference rules in (Li, Manyà, & Planes 2007) can be applied to $IS_1$ and $IS_2$ as they both contain a ternary clause.
Our algorithm for Max-SAT with the new method proposed in this section to compute lower bounds is described in Algorithm 1, called IncMax. Given a formula \( F \) and a ratio \( \alpha \), IncMax is called as \( \text{IncMax}(F, \#\text{clauses}, 0, 0, \alpha) \). The algorithm begins with handling some trivial cases (lines 3-4). Then it applies inference rules (like in (Li, Manyà, & Planes 2007)) to \( F \) (line 5). If \( UB \) does not reach \( \alpha \), the procedure identifyISS exploits unit propagation to identify disjoint inconsistent subformulas in \( F \), and \( LB \) is computed (lines 6-8). Otherwise, \( LB \) is computed in an incremental way (lines 9-16): Remove the clauses already in \( ISS \) in order to avoid re-computation; exploits unit propagation to compute \( LB \) in remaining clauses; and then reinsert the clauses in \( ISS \) into \( F \). If \( LB \) reaches \( UB \) the algorithm backtracks, otherwise a variable \( x \) is selected to instantiate (lines 17-19). If \( x \) is contained by an inconsistent subformula \( S \), unit propagation is exploited to reduce it (lines 20-25). Finally, the instance is solved by two recursive calls for the two sub-instances \( F[x] \) and \( F[x] \) (lines 26-27). Note that when \( \alpha = 1 \), IncMax is the same as the algorithm proposed by (Li, Manyà, & Planes 2007).

The correctness of IncMax is guaranteed by Proposition 1 and 2. As for complexity, IncMax only uses extra linear space to store \( ISS \) compared with the algorithm proposed by (Li, Manyà, & Planes 2007).

### Learning from Past

So far, our algorithm has not incorporated failed literal detection (FL). This is because if an inconsistent subformula \( IS \) is identified by FL, we cannot guarantee that the unsatisfiable core can be extracted by unit propagation (UP) after instantiating variables in it. Although, of course, it can be extracted by FL, it seems costly because FL is more time-consuming than UP. However, FL is very powerful for pruning branches on some instances, especially on Max-3-SAT and Max-CUT ones (Li, Manyà, & Planes 2006). So we extend our algorithm to a new version, called IncMax*, which exploits FL after UP to detect additional inconsistent subformulas, but such inconsistencies will not be inherited by the subsequent nodes.

Moreover, we found that, from empirical investigation, FL can rarely lead to backtrackings, e.g., only 5% of executions of FL making \( LB \) exceed \( UB \), on some instances. Thus, IncMax* enables the search-node to learn information about the effect FL has on the previous nodes. More specifically, at each node, FL is executed only if

\[
\#FL \leq \#\text{SAMPLE} \quad \text{or} \quad \frac{\#\text{FAIL} \times LB}{\#FL \times UB} \geq \beta \tag{1}
\]

where \#FL denotes the current total times of executions of FL, and \#FAIL denotes the current number of failures (backtrackings) led by FL. The condition (1) makes sure that for the first \#SAMPLE times, FL is always executed. After then, whether FL is executed depends on what percentage of the former executions have led to backtrackings and how close is \( LB \) to \( UB \). Note that, for an input formula, parameters \( \beta \) and \#SAMPLE are automatic determined according to the number of variables \( n \), the number of clauses \( m \), and the maximum clause length \( k \). Since FL seems powerful when \( k \geq 3 \), we set \( \beta = 0.2 \) and \( \beta = 0.3 \) when \( k \geq 3 \) and \( k \leq 2 \) respectively. For \#SAMPLE, ideally, it should be proportional to the size of the entire search tree. However, it is difficult to determine such size before search. So we set \#SAMPLE to be \( n \times m / 10 \) and \( n \times m / 100 \) when \( k \geq 3 \) and \( k \leq 2 \) respectively.

### Experimental Investigation

We have implemented IncMax and IncMax* in two new Max-SAT solvers called incmaxsat and incmaxsat+fl, which were accomplished by extending the Max-SAT solvers maxsat and maxsat+fl respectively. We conducted an experimental investigation to compare our solvers with the following state-of-the-art solvers.

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**Algorithm 1: IncMax**

```
1 IncMax(F, UB, LB, ISS, \alpha)
2 Input: A CNF formula \( F \), an upper bound \( UB \), a lower bound \( LB \), a set of disjoint inconsistent subformulas \( ISS \) and a constant ratio \( \alpha \).
3 Output: The minimum number of unsatisfied clauses in \( F \).
4 begin
5 if \( F = \emptyset \) or \( F \) only contains empty clauses then
6 return \#Empty(F)
7 F \leftarrow \text{applyRules}(F)
8 if \( LB < UB \times \alpha \) then
9 C \leftarrow \text{identifyISS}(F)
10 LB \leftarrow \#Empty(F) + \#C
11 else if \( LB < UB \) then
12 foreach inconsistent subformula \( S \in ISS \) do
13 remove clauses of \( S \) from \( F \)
14 C \leftarrow \text{identifyISS}(F)
15 LB \leftarrow \#Empty(F) + \#ISS + \#C
16 foreach inconsistent subformula \( S \in ISS \) do
17 remove clauses of \( S \) into \( F \)
18 ISS \leftarrow ISS \cup C
19 if \( LB > UB \) then
20 return \infty
21 x \leftarrow \text{selectVariable}(F)
22 if \( LB \geq UB \times \alpha \) then
23 foreach inconsistent subformulas \( S \in ISS \) do
24 if \( S \) contains \( x \) then
25 remove \( S \) from \( ISS \)
26 \( S' \leftarrow \text{UP}(S) \)
27 ISS \leftarrow ISS \cup \{S'\}
28 UB \leftarrow \min(UB, \text{IncMax}([\bar{x}], UB, LB, ISS, \alpha))
29 return \min(UB, \text{IncMax}(F[x], UB, LB, ISS, \alpha))
30 end
```

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\(^3\)In our solver, as well as maxsat, some inference rules are applied after detecting inconsistent subformulas. However, for simplicity, the pseudo codes do not reflect this case.
• maxsat\textsuperscript{6} (Li, Manyà, & Planes 2007): a branch and bound solver which exploits unit propagation and some sophisticated inference rules to compute lower bounds.
• maxsat+fl\textsuperscript{7} (Li, Manyà, & Planes 2006; 2007): a solver extending maxsat by incorporating failed literal detection. We used the evaluation version which turned out to be the best performing solver on unweighted instances in the Max-SAT Evaluation 2006 and 2007.
• maxsat14icss (Darras et al. 2007): a solver, built on maxsat+fl, using a heuristic to memorize inconsistencies.
• Toolbar3.0\textsuperscript{8} (Larrosa & Heras 2006; Larrosa, Heras, & de Givry 2008): a Max-SAT solver which exploits resolution rules and weighted CSP solving techniques extensively.
• MiniMaxSat\textsuperscript{9} (Heras, Larrosa, & Oliveras 2008): a new solver which incorporates many current SAT techniques, like clause learning and two-watched literal scheme.

All the benchmarks we used are from the Max-SAT Evaluation 2006 and 2007. Executions are run on a 1.7 Ghz Pentium 4 computer with 512 Mb of RAM under Linux.

In the first experiment we evaluated, on random 100-variable Max-2-SAT instances, the lower bound computation method by “learning from one’s own parent”. The number of clauses ranged from 300 to 1300. Figure 1 and Figure 2 show that the new method obtains better lower bounds than maxsat. As the number of clauses increases, the power of the new method is clearer.

As for random 70-variable Max-3-SAT instances, the lower bounds obtained by our method seem close to maxsat (see Figure 4). However, since the lower bound computing procedure is more efficient due to sparing re-computation time, incmaxsat still outperforms maxsat (see Figure 3).

On random 140-variable Max-2-SAT instances, we compared incmaxsat+fl with all other solvers. The results (see Figure 5) demonstrate that incmaxsat+fl is consistently better than others, especially when the number of clauses is large. On 80-variable Max-3-SAT instances, incmaxsat+fl is as good as maxsat+fl, and faster than others (Figure 6). As for structured instances (Table 1), incmaxsat+fl is the best performing one on most classes of benchmarks.

\textsuperscript{6}The source code is downloadable from http://web.udl.es/usuaris/m4372594/soft/maxsatz.tgz.
\textsuperscript{7}http://web.udl.es/usuaris/m4372594/soft/maxsatz-evaluation
\textsuperscript{8}http://mulcyber.toulouse.inra.fr/
\textsuperscript{9}http://www.lsi.upc.edu/~fheras/docs/m.tar.gz
Benchmarks | #Instances | Toolbar3.0 | MiniMaxSat | maxsatz+ff | maxsatz14icss | incmaxsatz+fl |
--- | --- | --- | --- | --- | --- | --- |
brock | 12 | 124.33(12) | 66.51(12) | 18.47(12) | 17.87(12) | **15.37(12)** |
c-fat | 7 | 209.59(5) | 0.23(5) | 0.11(5) | 0.11(5) | **0.09(5)** |
hamming | 6 | 1102.95(3) | 834.74(3) | 243.19(3) | 209.37(3) | **218.07(3)** |
johnson | 4 | 238.46(3) | 254.95(3) | 62.29(3) | 58.29(3) | **49.93(3)** |
keller | 2 | 27.67(2) | 18.92(2) | 8.59(2) | 6.84(2) | **6.02(2)** |
p-hat | 12 | 177.06(12) | 77.39(12) | 21.79(12) | 21.32(12) | **19.80(12)** |
san | 11 | 141.07(7) | 980.18(10) | 366.78(11) | 371.16(11) | **332.53(11)** |
sanr | 4 | 591.05(4) | 354.76(4) | 97.10(4) | 92.44(4) | **77.17(4)** |
random-maxcut | 40 | 98.89(40) | 23.28(40) | 8.13(40) | 7.60(40) | **7.90(40)** |
spinglass | 5 | 10.67(2) | 2.19(3) | 63.45(3) | 38.40(3) | **43.49(3)** |

Table 1: Mean time on structured benchmarks. The number of solved instances (within one hour) is displayed in brackets, and the best solver is highlighted in bold text.

Figure 6: Mean time for Max-3-SAT

Conclusions and Future Work

We have presented within-problem learning techniques, based on the similarity (Propositions 1 and 2) among the nodes of the search tree, for improving lower bound computation in Max-SAT solving. We implemented Max-SAT solvers, which can be considered as generalized versions of maxsatz. The experimental investigation demonstrates that the new solvers are very competitive.

As future work we plan to extend the within-problem techniques to weighted Max-SAT, partial Max-SAT and even other optimization problems. The techniques are promising because we believe that the similarity, among nodes of a search tree, commonly exists in many problems.

References


