Generating Parallel Execution Plans with a Partial-Order Planner*

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Abstract
Many real-world planning problems require generating plans that maximize the parallelism inherent in a problem. There are a number of partial-order planners that generate such plans; however, in most of these planners it is unclear under what conditions the resulting plans will be correct and whether the planner can even find a plan if one exists. This paper identifies the underlying assumptions about when a partial plan can be executed in parallel, defines the classes of parallel plans that can be generated by different partial-order planners, and describes the changes required to turn UCPOP into a parallel execution planner. In addition, we describe how this planner can be applied to the problem of query access planning, where parallel execution produces substantial reductions in overall execution time.

Introduction
There are a wide variety of problems that require generating parallel execution plans. Partial-order planners have been widely viewed as an effective approach to generating such plans. However, strictly speaking, a partially-ordered plan represents a set of possible totally-ordered plans. Just because two actions are unordered relative to one another does not imply that they can be executed in parallel. The semantics of a partially-ordered plan provide that the two actions can be executed in either order. Simultaneous execution requires that potential resource conflicts between unordered actions be made explicit and avoided.

There are numerous partial-order planners presented in the literature, including SIPE (Wilkins 1984), NONLIN (Tate 1976), SNLP (McAllester & Rosenblitt 1991), UCPOP (Penberthy & Weld 1992), TWNEAK (Chapman 1987), O-PLAN (Currie & Tate 1991), etc. Many of these planners have been used to produce parallel plans, but previously no one has precisely described the class of parallel plans they produce, identified their limitations, or described the assumptions made for the parallel plans that they do produce.

This paper focuses on the use of partial-order planning to generate parallel execution plans. First, we identify the conditions under which two unordered actions can be executed in parallel. The component missing from many planners is an explicit representation of resources. Second, assuming that the resource constraints have been made explicit, we identify the classes of parallel execution plans that can be generated using different partial-order planners. Third, we present an implementation of a parallel execution planner based on UCPOP. Fourth, we describe how this planner can be used to generate parallel query access plans. Fifth, we compare the use of a partial-order planner to other approaches to building parallel execution plans. Finally, we review the contributions of the paper and describe some directions for future research.

Executing Actions in Parallel
Classical planners assume that the execution of an action is indivisible and uninterruptible (Weld 1994). This is referred to as the atomic action assumption and stems from the fact that the STRIPS-style representation only models the preconditions and effects of all actions. This assumption would appear to make simultaneous execution impossible, since it is unclear from the action model whether any two actions can be executed simultaneously without interacting with one another. This section identifies the conditions under which it is possible to execute two actions in parallel.

The work on parallelizing execution of machine instructions (Tjaden & Flynn 1970) provides some insight on the types of dependencies that arise between actions. Tjaden and Flynn identify three types of dependencies that must be considered in parallelizing machine instructions: procedural, operational, and data. A procedural dependency occurs when one instruction explicitly addresses another instruction and therefore imposes an ordering between the instructions. An operational dependency occurs when there is a resource...
associated with an instruction that is potentially unavailable. A data dependency occurs when one instruction produces a result that is used by another instruction.

Similar dependencies arise in the parallelization of planning operations. A procedural dependency arises when one operation is explicitly ordered after another operation, which occurs in many of the hierarchical planners (Tate 1976; Wilkins 1984) (e.g., see the plot construct in SIPE). This type of constraint is captured by explicit ordering constraints between actions. A data dependency arises when the precondition of one operation depends on the effects of another operation. This type of dependency is captured by the operator representation and corresponding algorithms, which ensure that if two actions are unordered relative to one another, their preconditions and effects are consistent. Operational dependencies can occur when there are limited resources associated with an operation. This type of dependency is often ignored by planning systems.

Executing actions in parallel requires explicit handling of potential resource conflicts. If two actions are left unordered in a partial-order plan, they can be executed in either order. In order to execute them in parallel, we must ensure that there are no potential conflicts that occur during execution. Most conflicts will be resolved in the process of ensuring that the preconditions and effects are consistent. However, because of the limited representation, the type of conflict that is not typically handled in a partial-order planner is when two actions require the same reusable resource. This type of resource conflict is not typically captured by the preconditions and effects because at the start of execution the resource is available and when execution completes it is available.

Despite the problem of potential resource conflicts, a number of partial-order planners have allowed simultaneous execution. They do so by either assuming the actions are independent (Tate 1978), augmenting the action representation to avoid resource conflicts (Wilkins 1984), or requiring the user to explicitly represent the conflicts in the preconditions and effects of the operators (Currie & Tate 1991). The approach of simply assuming that the actions are independent could lead to unexpected resource conflicts. The approach of requiring the user to represent the conflicts in the preconditions and effects is both awkward and computationally more expensive, since it requires additional operators that explicitly allocate and deallocate resources. The most natural approach is to augment the action representation to describe the explicit resource needs of the different actions. This approach was proposed in SIPE (Wilkins 1984), where each operator can be annotated to explicitly state if something is a resource. In the next section we will assume that the resource constraints have been made explicit and in the following section we will describe our approach to representing and reasoning about resources.

Parallel Execution Plans

This section identifies the classes of parallel execution plans that can be generated by different planners, assuming that a domain is correctly axiomatized and explicitly represents the resource requirements of the operators. The different types of parallel execution plans can be broken down into several classes, ranging from plans with completely independent actions to those where the actions must be executed in parallel or must overlap in a particular manner in order for the plan to succeed. As the interactions in the plan increase in complexity, the corresponding planners require more sophisticated representations and reasoning in order to generate such plans. In this section we present four classes of parallel execution plans and identify the corresponding planners that can generate that class of plans. These classes are described in order from the most restrictive to the least restrictive.

Independent Actions

The most restricted type of parallel execution plans are those where all of the parallel actions are completely independent of one another.

Two actions are defined to be independent if and only if the effects of the two actions executed simultaneously are the union of the individual effects of the actions done in isolation.

Allen (1991) notes that various partial-order planners, such as NONLIN (Tate 1976), DEVISER (Vere 1983), and SIPE (Wilkins 1984), all “allow simultaneous action when the two actions are completely independent of each other.” While this statement is correct, it is a bit misleading since these planners can generate plans for a less restrictive class of parallel plans. As noted by Horz (1993), since some of the effects of an operator may be unnecessary with respect to the goal and preconditions of other operators, the fact that two operators are unordered in a plan generated by a partial-order planner does not imply that they are independent. A planner that can only generate plans with independent actions is UA (Minton, Bresina, & Drummond 1991), which imposes ordering constraints between any pair of unordered actions that could possibly interact.

Figure 1 illustrates a simple plan with two independent actions. The goal of the plan is to have the table painted red and the chair painted blue. Since the actions of painting the table and painting the chair are independent, they can be executed in parallel.

Independent Actions Relative to a Goal

In a variety of partial-order planners, such as SIPE (Wilkins 1984), SNLP (McAllester & Rosenblit 1991), and UCPOP (Penberthy & Weld 1992), the planners enforce the property that two actions can only remain unordered if there is no threat between them. A threat occurs when an operator could potentially delete a con-
Independent Subplans Relative to a Goal

Not all partial-order planners enforce the property that two actions can remain unordered only if there are no threats between them. In particular, those planners that implement some form of Chapman’s white knight (Chapman 1987) require only that there exist some operator that establishes a given precondition, but do not commit to which operator. More specifically, the white knight operation allows plans with the following conditions: There exists some operator \( op_1 \) that achieves a goal or precondition \( g \). There exists a second operator \( op_2 \) that possibly deletes \( g \). And there exists a third operator \( op_3 \) that follows \( op_2 \) and achieves \( g \). If we are interested in producing totally-ordered plans, then the white knight operator is not required for completeness. However, the use of the white knight operator allows a planner to generate a slightly more general class of parallel plans.

The planners in this class include TWEAK (Chapman 1987), NONLIN (Tate 1976), o-PLAN (Currie & Tate 1991), MP, and MPI (Kambhampati 1994). The class of parallel plans produced by these planners are those with independent subplans relative to a goal.

Two subplans are independent relative to a goal \( G \) if and only if, for all conditions that are relevant to achieving \( G \), the result of executing the subplans in either order is identical to the result of executing the actions simultaneously.

These planners are limited to this class of plans since, if there is a threat between any pair of actions, additional constraints are imposed on the plan to eliminate the threat.

Figure 2 shows an example plan where all of the actions are independent relative to the goal. This example differs from the independent action example in that the two painting actions each have a side-effect of painting the floor as well as the object. Thus, the paint table and paint chair operators are not independent since both operations also paint the floor different colors. However, since the color of the floor is irrelevant to the goal of getting the table painted red and the chair painted blue, the plan is still valid and could be generated by planners in this class.

Figure 3 shows an example plan with independent subplans relative to the goal (adapted from an example in (Kambhampati 1994)). In this example, before the table and chair can be painted red, they must be primed, and priming them has a side effect of painting the floor white. The final goal of the problem is to get the table, chair, and floor all painted red. Notice that the action of priming the chair interacts with painting the table, since they both change the color of the floor. Similarly, priming the table interacts with painting the chair. Despite these potential interactions, the floor will still be painted red at the end of the plan since the table and chair must be painted after they are primed. Solving this problem requires the white knight operation to produce the parallel plan since the plan does not state which painting operation will be used to achieve the final goal of making the floor red.

The implementation of the white knight, which allows a planner to generate this more general class of parallel plans, also makes it difficult to extend the operator language to efficiently handle more expressive constructs, such as conditional effects and universal quantification (Chapman 1987). These more expressive language constructs are often required for representing and solving real problems.

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1 Some planners, such as SNLP and earlier versions of UCPOP, defined a threat to include an operator that adds or deletes a relevant condition. This stronger definition of a threat is used to constrain the search space and would prevent some possible parallel plans from being generated.
Interacting Actions

The most general class of parallel plans are those where the parallel actions interact in some way. Two actions may need to be executed in parallel or two actions may need to overlap in a particular manner in order for the plan to succeed. For example, if the final goal was to get the chair blue, the table yellow, and the floor green and there was no green paint, we could paint the table and chair simultaneously. To handle these cases requires the introduction of an explicit representation of time, such as that provided in temporal planning systems (Allen et al. 1991; Penberthy 1993). However, in this paper we are interested in the more restricted case where we would like to execute actions in parallel to take advantage of the possible parallelism to reduce the total execution time, not because the solution requires parallelism to solve the problem.

Parallel Execution Planning in UCPOP

We used the ucPOP planner (Penberthy & Weld 1992; Barrett et al. 1993) to build a parallel execution planner. The analysis in the previous section showed that UCPOP can produce the class of plans with actions that are independent relative to a goal. For the specific application described in the next section, this restriction does not prevent the system from finding any solutions. The changes to UCPOP that were required were to add explicit resource definitions to the operators, to modify the planner to enforce the resource constraints, and to construct an evaluation functions to estimate the cost of the parallel plans.

The resource requirements of the operators are made explicit by augmenting each operator with a resource declaration. An example operator with a resource declaration is shown in Figure 4. This operator describes the action of moving data from one data source to another and declares the data source from which the data is being moved as a resource. The purpose of this declaration is to prevent one operator from being executed in parallel with another operator that requires the same database.

In order to avoid resource conflicts, we modified the planner to ensure that if two operators require the same resource, then they are not left unordered relative to one another. In SIPE this is done with a critic that checks for resource conflicts and then imposes ordering constraints when conflicts are found. In UCPOP, we added a check to the planner such that every time a new action is added to the plan, the planner checks for potential resource conflicts with any other operator that could be executed in parallel. Any conflicts discovered are added to the list of threats that must be removed before the plan is considered complete. Using the search control facility in UCPOP, these conflicts can be resolved immediately or delayed until later in the planning process.

Since efficiency is the primary motivation for generating parallel plans, we constructed an evaluation function that can be used to find plans with low overall execution time. Since this evaluation function underestimates the cost of the parallel plan, the planner can use a best-first search to find the optimal plan. This evaluation function takes into account that the cost of executing two actions in parallel will be the maximum and not the sum of the costs. The space of parallel execution plans may be quite large, so domain-specific control knowledge may be necessary to search this space efficiently.

The evaluation function to determine the execution time of a parallel execution plan is implemented using a depth-first search. The search starts at the goal node and recursively assigns a cost to each node in the plan. This cost represents the total cost of execution up to and including the action at the given node. The cost is calculated by adding the cost of the action at the node to the maximum cost of all the immediately prior nodes. Once the cost of the plan up to a node has been computed, we store this value so it will only need to be calculated once. Since each node (n) and each edge (e) in the graph is visited only once, the complexity of evaluating the plan cost is O(max(n,e)).
Parallel Query Access Plans

There are two general characteristics of a domain where the use of a partial-order parallel-execution planner will be useful and effective. First, it is applicable to those domains where the actions could be executed serially, but the overall execution time can be reduced by executing some of the actions in parallel. Second, it will be most useful in those domains where the choice of the operations determines or limits the overall execution time of the plan. As such, the plan generation and scheduling cannot be done independently since this would potentially result in highly suboptimal plans.

We applied the parallel execution planner to a query access planning problem that involves multiple distributed information sources (Knoblock & Arens 1994). In this domain, a plan is produced that specifies how to retrieve and generate the requested set of data. This requires selecting sources for the data, determining what operations need to be performed on the data, and deciding on the order in which to execute the operations. The planner must take into account the cost of accessing the different information sources, the cost of retrieving intermediate results, and the cost of combining these intermediate results to produce the final results. A partial-order parallel-execution planner is ideally suited for this problem since the parallelization is for efficiency purposes, and there are many possible plans for retrieving the same data and the choice of plans is crucial in determining the overall efficiency.

Figure 5 shows an example parallel query-access plan. The three basic query access planning operations used in this plan are move-data, join, and retrieve. The move-data operation moves a set of data from one information source to another. The join operation combines two sets of data into a combined set using the given join relations. The retrieve operation specifies the data that is to be retrieved from a particular information source.

The domain and planner described here are fully implemented and serve as an integral part of an information retrieval agent. We have also extended UCPOP to perform execution, and to do so in parallel. The system is implemented and runs in Lucid Common Lisp on SUN and HP workstations. To provide a sense for the potential speed-up of this approach we ran a sample query that involved queries to two different databases. Without parallelization, the system generated a plan with six operators in 0.82 CPU seconds and then executed the plan in 101.8 seconds of elapsed time. With parallelization, it generated the plan in 1.3 CPU seconds and executed the plan in 62.4 seconds of elapsed time, a 39 percent reduction in execution time.

Related Work

An alternative approach to addressing the problem of simultaneous execution is provided by work on temporal planning (Allen et al. 1991; Penberthy 1993). A temporal planner can handle the general problem of simultaneous parallel execution, but this general solution has a cost, since just testing the satisfiability of a set of assertions is NP-hard (Vilain, Kautz, & van Beek 1989). The capabilities of a full-fledged temporal planner are necessary only if we need to explicitly reason about the interaction between parallel actions. In this paper we focus on the simpler problem of non-interacting simultaneous execution, which does not require a full-blown temporal reasoner to handle. In fact, partial-order planners appear to be well suited for problems in this class.

Another approach to this problem is to generate totally-ordered plans and then convert each plan into a partially-ordered plan (Veloso, Perez, & Carbonell 1990: Regnier & Fade 1991). The problem with this approach is that the particular choice of the totally-ordered plan determines the parallel execution plan. As such, in order to consider the space of parallel execution plans requires searching through the space of totally-ordered plans. Since a single partially-ordered plan often corresponds to a number of totally-ordered plans, it will be harder to efficiently search the space of parallel execution plans.

Recently, Backstrom (1993) showed that the general problem of finding an optimal parallel execution plan
The idea of using partial-order planning to generate parallel execution plans has been around since the early days of planning. What we have done in this paper is to explicate the underlying assumptions and situations where parallel execution is possible, characterize the differences in the plans produced by various planning algorithms, and identify the changes required to use UCPOP as a parallel execution planner. We have also shown that these ideas apply directly to the problem of generating parallel query access plans.

In future work we plan to tightly integrate the planning and execution components. This would allow the system to dynamically replan actions that fail, while continuing to execute other actions that are already in progress. In addition, we plan explore the problem of how to efficiently search the space of parallel execution plans. First, we will consider domain-independent search strategies that produce the highest quality solution that can be found within the time allotted. Second, we will exploit domain-specific knowledge to both restrict the search space and guide the search.

Discussion

The idea of using partial-order planning to generate parallel execution plans has been around since the early days of planning. What we have done in this paper is to explicate the underlying assumptions and situations where parallel execution is possible, characterize the differences in the plans produced by various planning algorithms, and identify the changes required to use UCPOP as a parallel execution planner. We have also shown that these ideas apply directly to the problem of generating parallel query access plans.

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